

DNSSEC

Stand der Einführung, Probleme und Entwicklungen

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Programm

- Warum DNSSEC?
- Wie funktioniert DNSSEC?
- Alternativen?
- Was heißt es, DNSSEC einzuführen?

Fragen sind willkommen!

Danke an Olaf Kolkman, Phil Regnauld, ... für die Erlaubnis, aus ihren Slides zu zitieren.

Warum überhaupt DNSSEC?

- Was macht das DNS so wichtig?
- Was sind die Angriffsszenarien auf das DNS?
- Der Kaminsky Attack.

DNS: Was ist das?

- Global, distributed Database
- Input: Domain name
- Output: Resource Record **Sets**
 - A, AAAA IP addresses
 - MX Mail routing
 - CNAME Aliasing
 - NS Delegation
 - PTR, NAPTR, SRV,
 - RRSIG, DS, NSEC, NSEC3
- Transport: mainly UDP
- Lots of caching

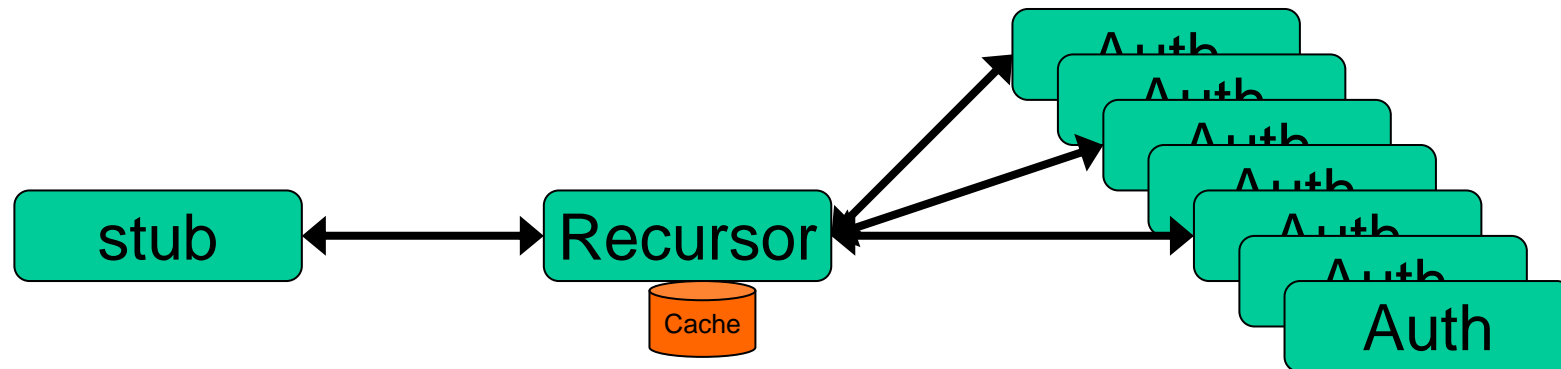


DNS: Warum ist es wichtig?

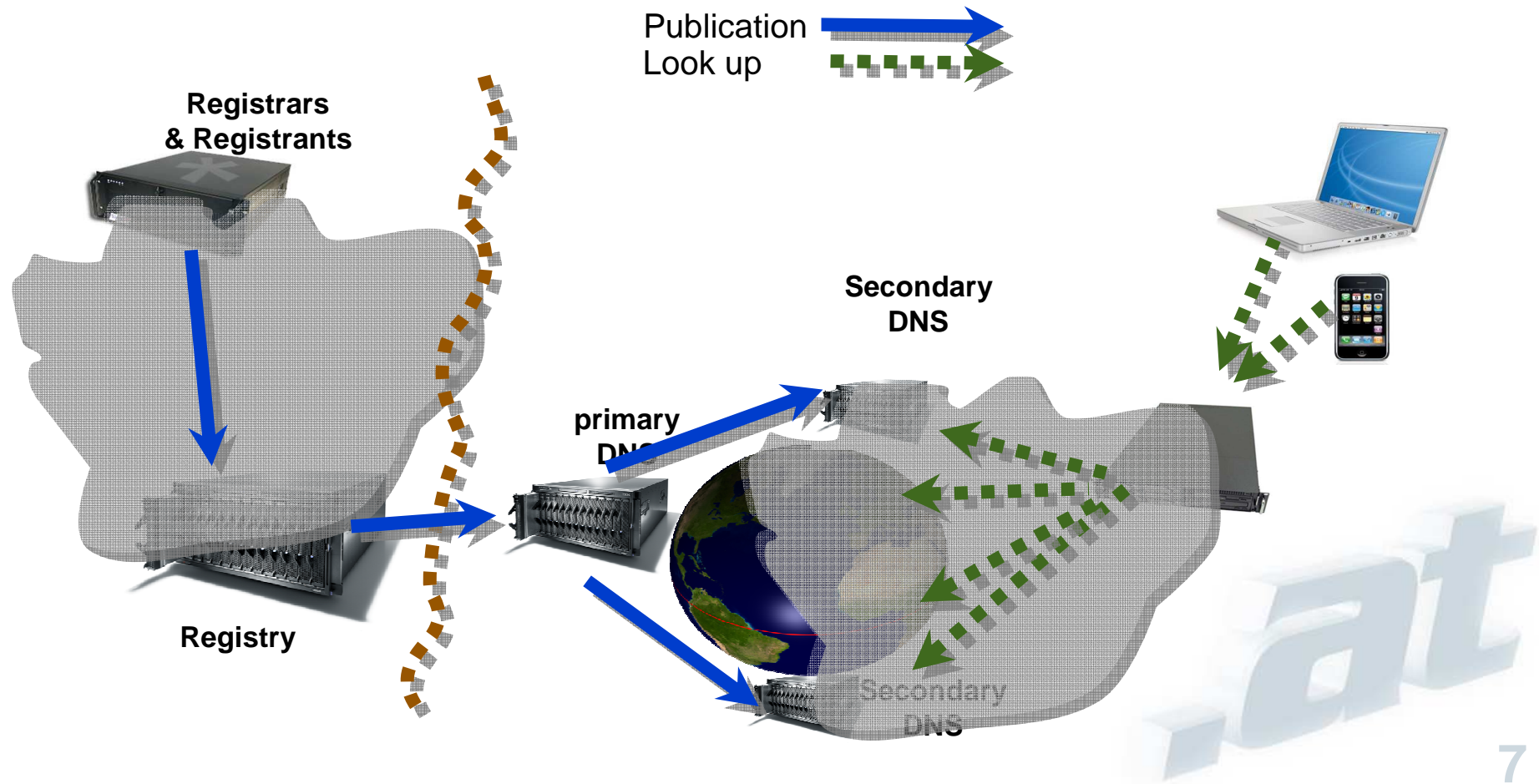
- DoS
- Man-in-the-middle almost *everything*
 - Phishing
 - Email hijacking
- Password reset emails
- Software Updates
- SSL and PKI for the rescue?
 - How do users react to X.509 errors?
 - CA email-loop
 - CA whois lookup
- Für den Enduser ist „DNS Kaputt“ nicht von „Internet ist kaputt“ unterscheidbar

Die Mitspieler

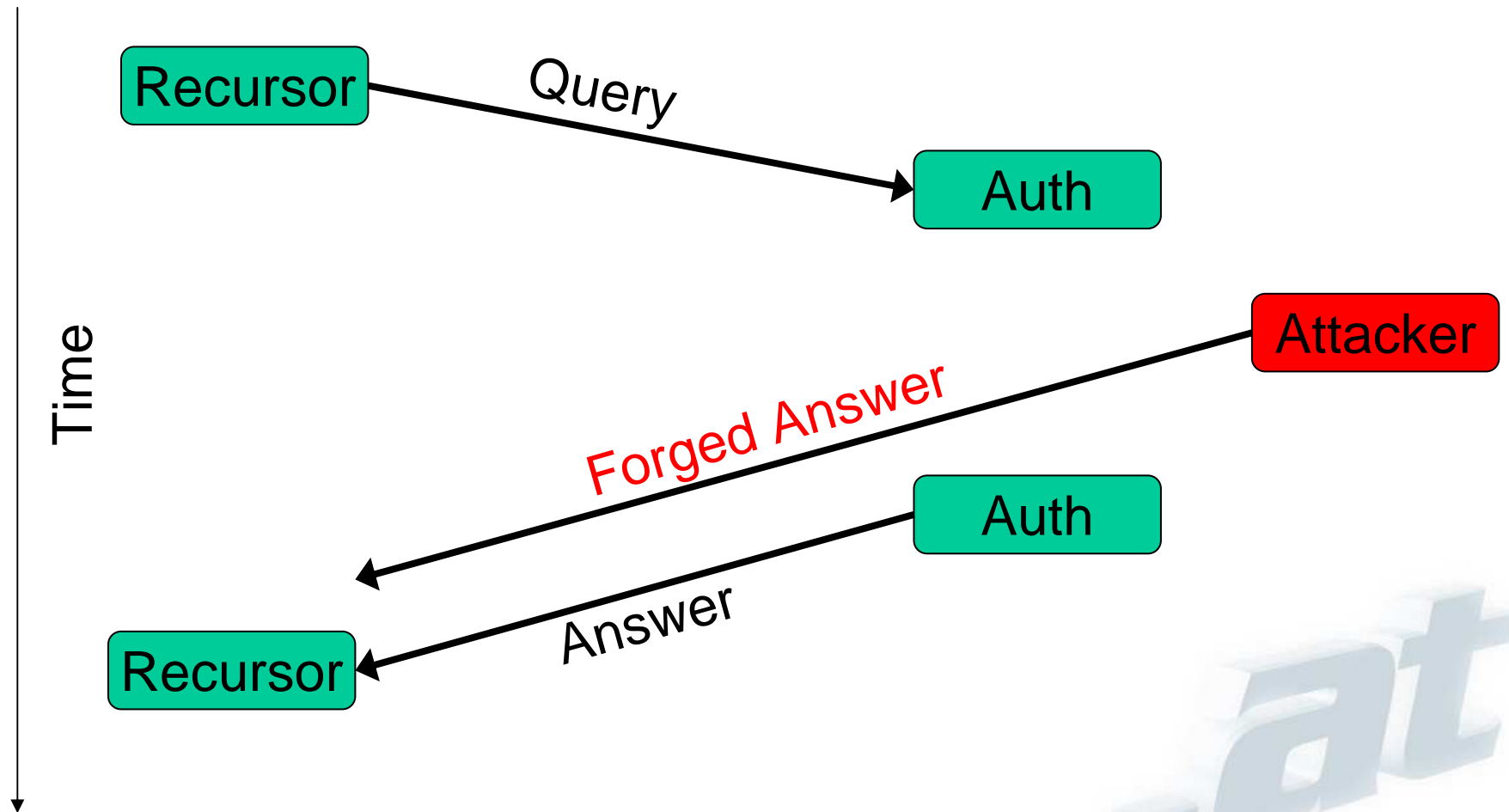
- Stub Resolvers
- Recursive Nameservers
- Authoritative Nameservers



Data flow through the DNS



DNS Spoofing



Gefahrenanalyse

- Cache poisoning
 - Off-path attacks
 - On-path attacks
- Name chaining
- Falsche Antwort durch den Recursor
 - Sitefinder
 - NXdomain monetizing

- Siehe auch RFC 3833



Cache Poisoning (off-path)

- Das ist nichts neues.
 - Kashpureff
 - Triviale Query-ID
 - Parallele Anfragen
- Berechnungen zur Erfolgswahrscheinlichkeit in RFC 5452



Pre-Kaminsky

- An attack needs to match
 - Question section
 - The ID field
 - IP address of the nameserver queried
 - IP address / port from which the query was sent
- How often can an attack take place?
 - Each query from a recursor starts a race.
 - Forcing a query helps the attacker
 - The cache limits attacks to once per Time-To-Live for the same query



Attacking www.example.org

;; QUESTION SECTION:

;345678.example.org. IN A

;; ANSWER SECTION:

345678.example.org. 3600 IN A 192.0.2.1

;; AUTHORITY SECTION:

example.org. 100000 IN NS ns1.evil.net.

example.org. 100000 IN NS ns2.evil.net.

Source: IETF namedroppers list. (P. Koch, T. Finch)

Oder ...

;; QUESTION SECTION:

;345678.www.example.org. A

;; AUTHORITY SECTION:

www.example.org. NS ns1.evil.net.

www.example.org. NS ns2.evil.net.



... oder ...

```
;; QUESTION SECTION:  
;345678.example.org.
```

```
IN A
```

```
;; ANSWER SECTION:  
345678.example.org.  
www.example.org.
```

```
CNAME www.example.org.  
A 192.0.2.80 ; evil
```



Wie funktioniert DNSSEC?

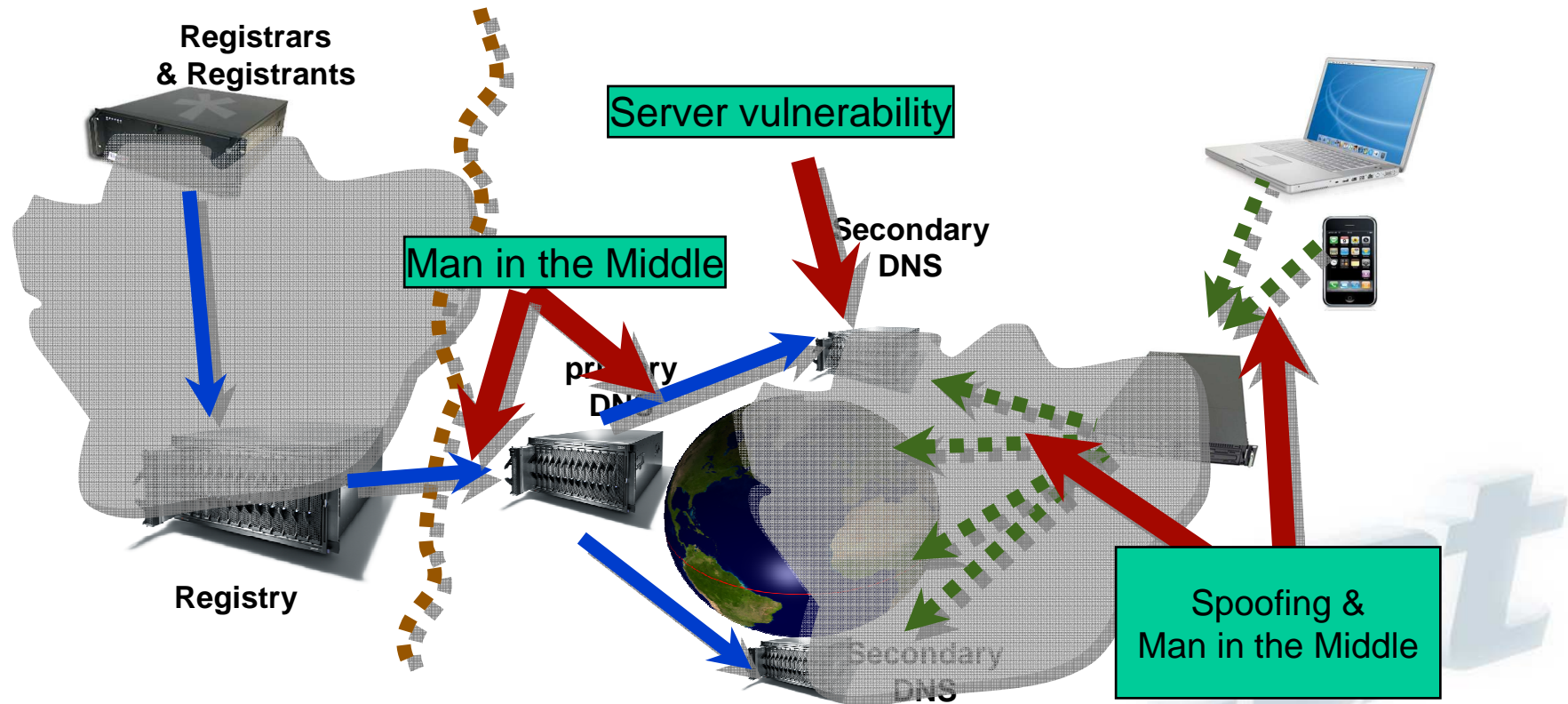
- Was soll DNSSEC leisten
- Die neuen Resource records
- Secure Delegations
- Key Management / Key Rollover
- NSEC3



DNSSEC Specs

- Details zum Nachlesen:
 - RFC4033, “DNS Security Introduction and Requirements”
 - RFC4034, “Resource Records for the DNS Security Extensions”
 - RFC4035, “Protocol Modifications for the DNS Security Extensions”
 - RFC5011, “Automated Updates of DNS Security (DNSSEC) Trust Anchors“
 - RFC5155, “DNS Security (DNSSEC) Hashed Authenticated Denial of Existence”

Where are the vulnerable points?



Welche Security?

- Confidentiality
 - Kann wer mitlesen?
- Integrity
 - Stimmt das, was ich bekommen habe?
- Availability
 - Bekomme ich überhaupt eine Antwort?

DNSSEC betrifft ausschließlich „Integrity“!

Grundidee

- Kompatible Erweiterung des DNS
- Public Key Kryptografie Signaturen innerhalb der DNS Antworten
- Schutz der Daten, nicht Schutz des Transports
- Delegationshierarchie des DNS wird auch zur Trust-Hierarchie



New Resource Records

- Three Public key crypto related RRs
 - RRSIG Signature over RRset made using private key
 - DNSKEY Public key, needed for verifying a RRSIG
 - DS Delegation Signer; 'Pointer' for building chains of authentication

- Two RR for internal consistency
 - NSEC Indicates which name is the next one in the zone and which typecodes are available for the current name.

 - NSEC3 NSEC++



RRSIG – Signature

- 16 bits - type covered
- 8 bits - algorithm
- 8 bits - nr. labels covered
- 32 bits - original TTL
- 32 bit - signature expiration
- 32 bit - signature inception
- 16 bit - key tag
- signer's name

RRSIGs gelten nicht ewig!

```
nlnetlabs.nl. 3600 IN RRSIG A 5 2 3600 (
  20050611144523 20050511144523 3112 nlnetlabs.nl.
  VJ+8ijXvbrTLeoAiEk/qMrdudRnYZM1VlqhN
  vhYuAcYKe2X/jqYfMfjfsUrmhPo+0/GOZjW
  66DJubZPmNSYXw== )
```

DNSKEY – Public Key

- 16 bits: FLAGS
- 8 bits: protocol
- 8 bits: algorithm
- $N \times 32$ bits: public key

```
nlnetlabs.nl. 3600 IN DNSKEY 256 3 5 (  
  AQOvhvXXU61Pr8sCwELcqqq1g4JJ  
  CALG4C9EtraBKVd+vGIF/unwigfLOA  
  O3nHp/cgGrG6gJYe8OWKYNgq3kDChN)
```



Delegation Signer (DS)

- Delegation Signer (DS) RR indicates that:
 - delegated zone is digitally signed
 - indicated key is used for the delegated zone

- Parent is authoritative for the DS of the child's zone
 - Not for the NS record delegating the child's zone!
 - **DS should not** be in the child's zone

DS – Key of Subdomain

- 16 bits: key tag
- 8 bits: algorithm
- 8 bits: digest type
- 20 bytes: SHA-1 Digest

```
$ORIGIN nlnetlabs.nl.  
lab.nlnetlabs.nl. 3600 IN NS ns.lab.nlnetlabs.nl  
lab.nlnetlabs.nl. 3600 IN DS 3112 5 1 (  
    239af98b923c023371b52  
    1g23b92da12f42162b1a9  
    )
```


NSEC – Proof of non-existence

- FQDN: Next Name in Zone
- N*32 bit map: RRTypes present

```
www.nlnetlabs.nl. 3600 IN      NSEC z.nlnetlabs.nl. A RRSIG NSEC
```



NSEC Records

- NSEC RR provides proof of non-existence
- If the servers response is Name Error (NXDOMAIN):
 - One or more NSEC RRs indicate that the name or a wildcard expansion does not exist
- If the servers response is NOERROR:
 - And empty answer section
 - The NSEC proves that the QTYPE did not exist
- More than one NSEC may be required in response
 - Wildcards
- NSEC records are generated by tools
 - Tools also order the zone



NSEC Walk

- NSEC records allow for zone enumeration
- Providing privacy was not a requirement at the time
- Zone enumeration is a deployment barrier

- Solution is developed: NSEC3
 - RFC 5155
 - Complicated piece of protocol work
 - Hard to troubleshoot
 - Only to be used over Delegation Centric Zones

DNSSEC Queries

- DO
 - DNSSEC OK (EDNS0 OPT header) to indicate client support for DNSSEC options
 - EDNS0 is required for DNSSEC
- CD
 - “Don’t check signatures for me, just give me the raw DNSSEC records”



DNSSEC Answers

- SECURE Validated with key
 - AD – bit set in Packet
- INSECURE Validated but no key
- BOGUS Validation failed
- UNKNOWN ServFail etc



Key management

- To allow for key updates (“rollovers”), generate two keys:
 - Key Signing Key (KSK)
 - ◆ pointed to by parent zone (Secure Entry Point), in the form of DS (Delegation Signer)
 - ◆ used to sign the Zone Signing Key (ZSK)
 - Zone Signing Key (ZSK)
 - ◆ signed by the Key Signing Key
 - ◆ used to sign the zone data RRsets
- This decoupling allows for independent updating of the ZSK without having to update the KSK, and involve the parent.

Deployment Server-side

- Key management
 - Generate keys
 - Add DNSKEY records
- Sign zone
 - Signing & serving need not be performed on same machine
 - Signing system can be offline
- Make sure authoritative nameservers handle DNSSEC
- Communicate your keys to parent zone

Deployment Client-side

- Stub-Resolver speaks DNSSEC
 - Inefficient
 - Slow rollout
 - Upsides in User-Interface
- Recursor does DNSSEC Validation
 - Need a way to secure last hop
 - Huge multiplier possibilities
- Secure Entry Points?



Trust Anchors

- Irgendwem muss der Client vertrauen
 - Hardcoded (domain, DNSKEY) Paare in der resolver-config
 - Analog zu dem „root.hints“ File
- Optimal:
 - Root ist signiert, alle TLDs,
- Realität:
 - Es gibt signierte Inseln:



Status 2009



Source: <http://www.xelerance.com/dnssec/>

Zwischenlösungen

- Selber Trust Anchors zusammensuchen
- DNSSEC Lookaside Validation (DLV)
 - <fqdn>.dlv.isc.org
- Trust Anchor Registries:
 - IANA Interim Trust Anchor Repository
 - ◆ <https://itar.iana.org/>
 - RIPE NCC?
 - ◆ <http://www.ripe.net/ripe/tf/dnssec-key/>
- Private, signed roots



Alternativen?

- On-path attacker?
 - Keine chance.
- Off-path attacker:
 - Mehr Entropie
 - ◆ 0x20
 - ◆ EDNS-PING
 - Skepsis bzgl. Zusatzinfos in DNS Antwort
 - ◆ Never cache data from auth and additional section.
 - ◆ Explicitly query for NS and A records.
 - ◆ Be careful when overwriting the cache.
 - Dynamisches reagieren:
 - ◆ Fallback to TCP
 - ◆ Multiple Queries

